## Indian Institute of Information Technology Allahabad

## Data Structures

## Linked List

## .. and other linked structures

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## Lists

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- Array - one way to represent a list.
- Constant time access given index of an element


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- Constant time access given index of an element
- Problems with arrays
- Size of an array should be specified beforehand (at least while dynamically allocating memory).
- Deleting/Inserting an element requires shifting of elements.
- Wasted space.


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Polynomial: $x^{\wedge} 25+3 x^{\wedge} 7-4$
Store in an array poly[26] poly[i] contains coefficient of $x^{\wedge}$ i
$\operatorname{poly}[0]=-4, \operatorname{poly}[7]=3, \operatorname{poly}[25]=1$
poly[i] $=0$ for all $i!=0,7,25$
Can we avoid storing so many 0's?
Store ( $0,-4$ ), $(7,3),(25,1)$ instead.
How do we 'link' these pairs?

## Dynamic Memory Allocation: Review

typedef struct \{ int hiTemp; int loTemp; double precip;
\} WeatherData;

```
int main ( ) {
int numdays;
WeatherData *days;
scanf ("\%d", \&numdays) ;
days=(WeatherData *)malloc (sizeof(WeatherData)*numdays);
if (days == NULL) printf ("Insufficient memory\n");
free (days) ;

\section*{Self-Referential Structures}

A structure referencing itself - how?


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So, we need a pointer inside a structure that points to a structure of the same type.
```

struct list {
int data;
struct list *next;
};

```

\section*{Self-Referential Structures}
struct list \{
int data ;
struct list *next ;
\};

The pointer variable next is called a link.
Each structure is linked to a succeeding structure by next.

\section*{Pictorial Representation}

A structure of type struct list


The pointer variable next contains either
- an address of the location in memory of the successor list element - or the special value NULL defined as 0 .

NULL is used to denote the end of the list.

\section*{Pictorial Representation}
struct list \(a, b, c\); a.data \(=1 ;\) b.data \(=2 ;\) c.data \(=3\); a.next \(=\) b.next \(=\) c.next \(=\) NULL;


\section*{Pictorial Representation}
\[
\begin{aligned}
& \text { a.next }=\& b ; \\
& \text { b.next }=\& c ;
\end{aligned}
\]


What are the values of :
- a.data
-a.next->data
-a.next->next->data

\section*{Pictorial Representation}
\[
\begin{aligned}
& \text { a.next }=\& b ; \\
& \text { b.next }=\& c ;
\end{aligned}
\]


What are the values of :
- a.data
-a.next->data 2
-a.next->next->data 3

\section*{Linked Lists}

A singly linked list is a concrete data structure consisting of a sequence of nodes
Each node stores
- element
- link to the next node


\section*{Linear Linked Lists}
- A head pointer addresses the first element of the list.
- Each element points at a successor element.
- The last element has a link value NULL.


\section*{Header File: list.h}
\#include <stdio.h>
\#include <stdlib.h>
typedef char DATA;
struct list \{
DATA d;
struct list *next;
\};
typedef struct list ELEMENT;
typedef ELEMENT *LINK;

\section*{Storage Allocation}
```

LINK head;
head = (LINK) malloc (sizeof(ELEMENT));
head->d = 'n';
head->next = NULL;

```
creates a single element list.


\section*{Storage Allocation}
```

head->next = (LINK) malloc (sizeof(ELEMENT));
head->next->d = 'e';
head->next->next = NULL;

```

A second element is added.


\section*{Storage Allocation}
head->next->next \(=(\) LINK \()\) malloc \((\) sizeof(ELEMENT) \()\);
head->next->next->d = 'w';
head->next->next->next = NULL;

We have a 3-element list pointed to by head.
The list ends when next has the sentinel value NULL.


\section*{List Operations}
- How to initialize such a self referential structure (LIST),
- How to insert such a structure into the LIST,
- How to delete elements from it,
- How to search for an element init,
- How to print it,
- How to free the space occupied by the LIST?

\section*{Produce a list from a string (Recursive Version)}
\#include <stdio.h> \#include <stdlib.h> typedef char DATA; struct list \{

DATA d;
struct list *next; \};
typedef struct list ELEMENT; typedef ELEMENT *LINK;

LINK StrToList (char s[ ]) \{
LINK head ;
if ( \(\mathrm{s}[0]==\) ' \(\backslash 0\) ') return NULL;
else\{
head \(=(\) LINK \()\) malloc (sizeof(ELEMENT));
head->d = s[0];
head->next = StrToList (s+1);
return head;

\section*{Produce a list from a string (Iterative Version)}
LINK SToL (char s[ ]) \{
    LINK head = NULL, tail; int i;
    if ( \(s[0]!=10\) ') \(\{\)
        head = (LINK) malloc (sizeof(ELEMENT));
        head->d = s[0];
        tail = head;

        tail->next=(LINK)malloc(sizeof(ELEMENT));
        tail = tail->next; tail->d = s[i];
        \}
    tail->next = NULL;
    \}
return head;
\}

tail head


\section*{Inserting at the Head}


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1. Allocate a new node
new = malloc(sizeof(ELEMENT));

\section*{Inserting at the Head}

1. Allocate a new node
2. Insert new element
new = malloc(sizeof(ELEMENT)); New->d = 'W';

\section*{Inserting at the Head}


\section*{Inserting at the Head}

1. Allocate a new node
2. Insert new element
3. Make new node point to old head
4. Update head to point to new node

new \(=\) malloc(sizeof(ELEMENT));
New->d = 'W';
new-> next = head;
head = new;
head


\section*{Removing the Head}

1. Update head to point to next node in the list
2. Allow garbage collector to reclaim the former first node

\section*{Removing the Head}


\section*{Removing the Head}


\section*{Removing the Head}


\section*{Inserting at the Tail}


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1. Allocate a new node
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\section*{Inserting at the Tail}

1. Allocate a new node
2. Insert new element
\[
\begin{aligned}
& \text { new = malloc(sizeof(ELEMENT)); } \\
& \text { new->d = ' } Z \text { '; }
\end{aligned}
\]

\section*{Inserting at the Tail}

1. Allocate a new node
2. Insert new element
3. Have new node point to null

new = malloc(sizeof(ELEMENT)); new->d = ' \(Z\) '; new->next = NULL;

\section*{Inserting at the Tail}

1. Allocate a new node
2. Insert new element
3. Have new node point to null
4. Have old last node point to new node

new = malloc(sizeof(ELEMENT)); new->d = ' \(Z\) '; new->next = NULL; tail->next = new;

\section*{Inserting at the Tail}

1. Allocate a new node
2. Insert new element
3. Have new node point to null
4. Have old last node point to new node
5. Update tail to point to new node
new = malloc(sizeof(ELEMENT)); new->d = 'Z'; new->next = NULL; tail->next = new; tail = new;

\section*{Removing the Tail}
1. Bring ptr to the second last node


\section*{Removing the Tail}


\section*{Removing the Tail}
1. Bring ptr to the second last node
2. Make ptr->next equal to NULL
3. Free tail
4. Make ptr the new tail


\section*{Insertion into an ordered list}

Create a new node containing the data

new = malloc(sizeof(ELEMENT));
New->d = 'N';

\section*{Insertion into an ordered list}

Create a new node containing the data


Find the correct place in the list


\section*{Insertion into an ordered list}

Create a new node containing the data


\section*{Insertion into an ordered list}

Create a new node containing the data


\section*{Insertion Function}
```

\#include <stdio.h>
\#include <stdlib.h>
struct list {
int data;
struct list *next;
};
typedef struct list ELEMENT;
typedef ELEMENT * LINK;
LINK create_node(int val) {
LINK newp;
newp = (LINK) malloc (sizeof
(ELEMENT));
newp -> data = val;
return (newp);
}

```
    LINK insert (int value, LINK ptr)
\{ LINK newp, prev, first;
    newp = create_node(value);
    if (ptr == NULL || value <= ptr -> data)
    \{ // insert as new first node
    newp -> next = ptr;
        return newp; // return pointer to first node
    \}
    else \{ // not first one
        first = ptr; // remember start
        prev = ptr;
        ptr = ptr-> next; // second
        while (ptr != NULL \&\& value > ptr -> data)
        \{ prev = ptr; ptr = ptr -> next; \}
        prev -> next = newp; // link in
        newp -> next = ptr; //new node
    return first;

\section*{Deletion}

\section*{Steps:}
- Finding the data item in the list


\section*{Deletion}

\section*{Steps:}
- Finding the data item in the list, and
- Linking out this node


\section*{Deletion}

\section*{Steps:}
- Finding the data item in the list, and
- Linking out this node, and
- Freeing up this node as free space.


\section*{Deletion}

\section*{Steps:}
- Finding the data item in the list, and
- Linking out this node, and
- Freeing up this node as free space.


What will happen if we did the following?
free(ptr); prev->next = ptr->next;

\section*{Deletion Function}
// delete the item from ascending list LINK delete_item(int val, LINK ptr) \{ LINK prev, first;
```

first = ptr; //remember start
if (ptr == NULL) return NULL;
else if (val == ptr -> data) //first node
{
ptr = ptr -> next; //second node
first->next = NULL;
free(first); //free up node
return ptr; //second
}

```
else //check rest of list
\{ prev = ptr;
    ptr = ptr -> next;
    // find node to delete
    while (ptr != NULL \&\& val > ptr->data)
    \{ prev = ptr; ptr = ptr -> next; \}
    if (ptr == NULL || val != ptr->data) \{
        // NOT found in ascending list
        return first; //original
    \}
    else \{ //found, delete ptr node
        prev -> next = ptr -> next;
        ptr->next = NULL;
        free(ptr); //free node
        return first; //original

\section*{Linear Searching for a data element in a list}
```

int Search( LINK head, int element) {
LINK temp;
temp = head;
while (temp != NULL) {
if (temp -> data == element) return 1;
temp = temp -> next;
}
return 0;
}

```

\section*{Printing a list}
void print_list (LINK head) \{
LINK temp;
temp = head;
while(temp!=NULL) \{
if(temp->next ==NULL)
printf("\%d. END OF LIST \(\ln "\), temp->data);
else printf("\%d -> ", temp->data);
temp = temp->next;
\}
\}

\section*{Printing a list backwards}

- How can you print backwards when the links are in forward direction?
- Can you apply recursion?

\section*{Printing a list backwards - recursively}
void PrintArray(LINK head) \{
if(head \(->\) next \(==\) NULL) \{ //boundary condition to stop recursion printf(" \%d, ",head -> data); return;
\}
PrintArray(head -> next);
printf(" \%d,",head -> data); return;
\}
//calling function recursively
//Printing current element

\section*{Freeing a list}
- What will happen if we free the first node of the list without placing a pointer on the second?

\section*{Freeing a list}
- What will happen if we free the first node of the list without placing a pointer on the second?
- In each iteration temp1 points at the head of the list and temp2 points at the second node.
```

void Free(ELEMENT *head) {
ELEMENT *temp1, *temp2;
temp1 = head;
while(temp1 != NULL) {
temp2 = temp1 -> next;
temp1->next = NULL;
free(temp1);
temp1 = temp2;
}
}

```

\section*{Counting the number of nodes in a list}
```

RECURSIVE APPROACH
int count (LINK head) {
if (head == NULL) return 0;
return 1+count(head->next);
}

```
```

ITERATIVE APPROACH

```
ITERATIVE APPROACH
int count (LINK head) {
int count (LINK head) {
    int cnt = 0;
    int cnt = 0;
    for (; head != NULL; head=head->next)
    for (; head != NULL; head=head->next)
    ++cnt;
    ++cnt;
    return cnt;
    return cnt;
}
```


## Concatenate two Lists

void concatenate (LINK ahead, LINK bhead) \{
if (ahead->next == NULL)
ahead->next = bhead;
else
concatenate (ahead->next, bhead);
\}
... And "Other" linked structures

- Like Trees, Sparse Matrices and Graphs

> Binary Tree
> (Height = 3)


Sparse Matrix
$\left[\begin{array}{llllll}0 & 9 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 5 \\ 6 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 3 & 0\end{array}\right]$

Graph
(Number of Vertices $=10$ ) (Number of Edges $=11$ )

... And "Other" linked structures: Binary Trees

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... And "Other"' linked structures: Sparse matrices
For the sparse matrix below:

$$
\left[\begin{array}{llllll}
0 & 9 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 5 \\
6 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 3 & 0
\end{array}\right]
$$

Storage as a 2-D Array: int M[5][6];

Storage required for 30 elements (with only 4 non-zero entries)

$$
=30 \text { * sizeof(int) = 120Bytes (For integers of size } 4 \text { Bytes) }
$$

... And "Other"' linked structures: Sparse matrices
Storage as a list of Tripples: (row, column, data)

$$
\left[\begin{array}{llllll}
0 & 9 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 5 \\
6 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 3 & 0
\end{array}\right]
$$

struct ripple \{ int row, column, data; struct ripple *next;
\}
head


## ... And "Other" linked structures: Graphs

## Adjacency Matrix Representation:

- Matrix location ( $\mathrm{i}, \mathrm{j}$ ) indicates an edge between vertices "i" and "j"
$\left[\begin{array}{llllll}0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0\end{array}\right]$


Storage as a 2-D Array: int G[6][6];

Storage required for 36 elements (with 6 vertices 5 edges)

$$
=36 \text { * sizeof(int) = 144Bytes (For integers of size } 4 \text { Bytes) }
$$

... And "Other" linked structures: Graphs

## Adjacency List Representation:

- Each vertex's neighbours are maintained in a linked list

struct vertex\{
int id;
struct vertex *next_adj;

Storage required $=(|\mathrm{V}|+$ sum of degree $) *$ sizeof( structure $)$

## Practice Problems

1. Concatenate two lists (iteratively)
2. Reverse a list
3. Delete the maximum element from a list
4. Rotate the list by k positions counter-clockwise
5. Write functions to create, insert, delete, display, search a sparse matrix

For each of the above, first create the linked list by reading in integers from the keyboard and inserting one by one to an empty list

## Acknowledgement

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## Thank You

