

Module 1 - Distributed System Architectures & Models

Architecture

- Distributed systems tend to be very complex.
- It is critical to properly organize these systems to manage the complexity.
- The organization of a distributed system is primarily about defining the software components that constitute the system.
 - A component is a modular unit with well-defined required and provided interfaces.

Architecture (2)

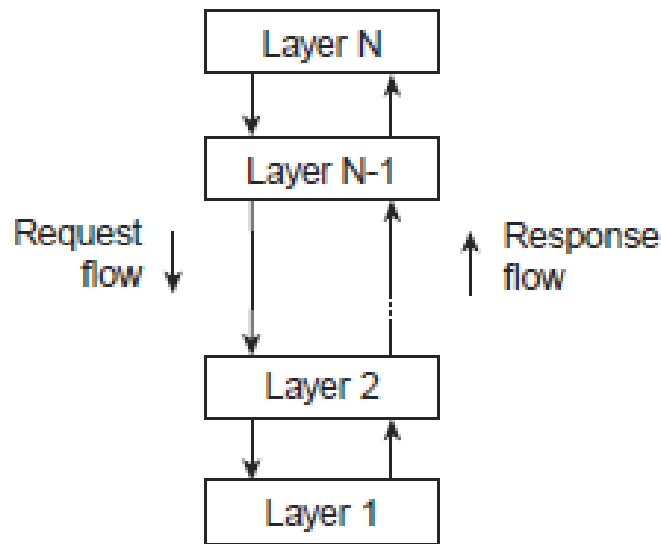
- Software Architecture:
 - Tells us how software components should be organized and how they should interact.

- System Architecture:
 - Instantiation of a software architecture on real machines.
 - Functions of each component are defined
 - Interrelationships and interactions between components are defined

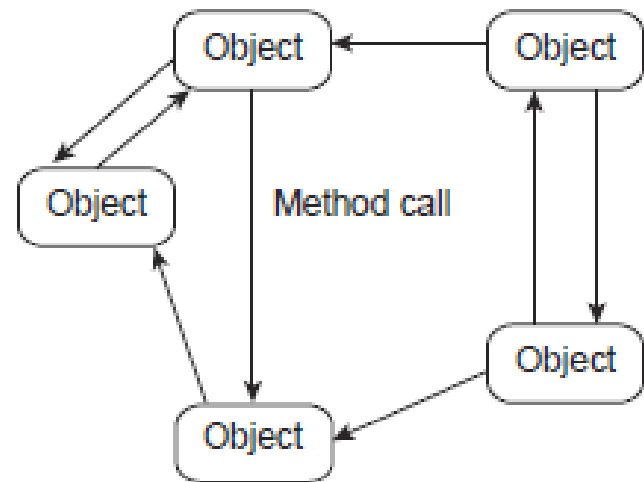
Architectural styles

Basic idea

Organize into **logically different** components, and distribute those components over the various machines.



(a)



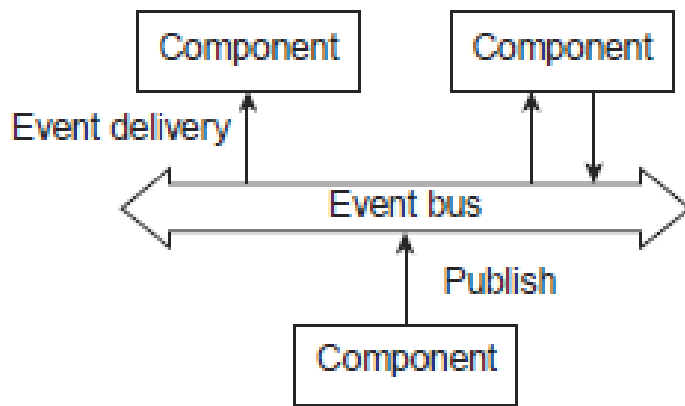
(b)

- (a) Layered style is used for client-server system
- (b) Object-based style for distributed object systems.

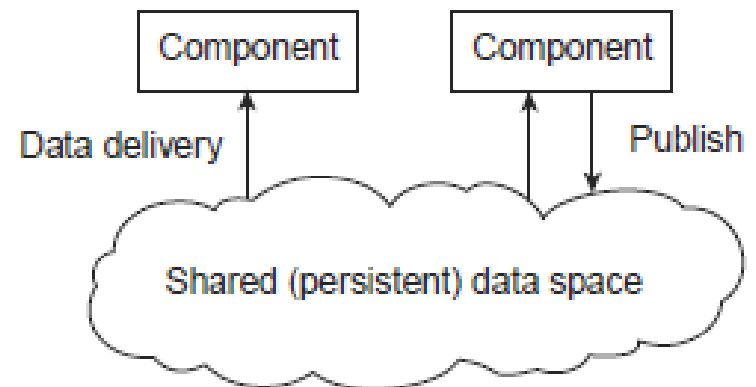
Architectural Styles

Observation

Decoupling processes in **space** (“anonymous”) and also **time** (“asynchronous”) has led to alternative styles.



(a)



(b)

(a) Publish/subscribe [decoupled in **space**]

(b) Shared dataspace [decoupled in **space** and **time**]

System Architectures

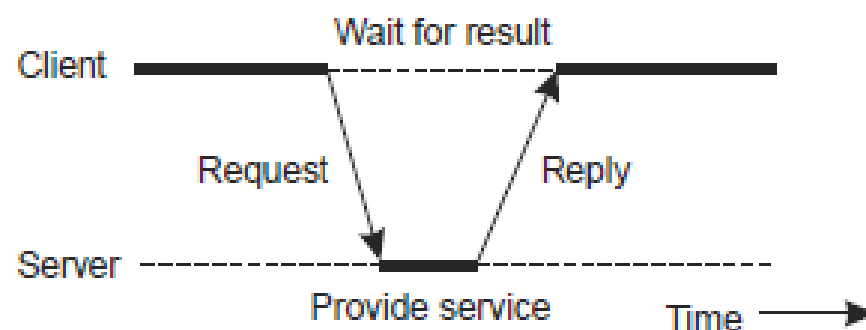
- Centralized architectures
 - Client-server
 - Multiple-client/single-server
 - Multiple-client/multiple-servers
 - Multitier systems
- Decentralized architectures

Centralized Architectures

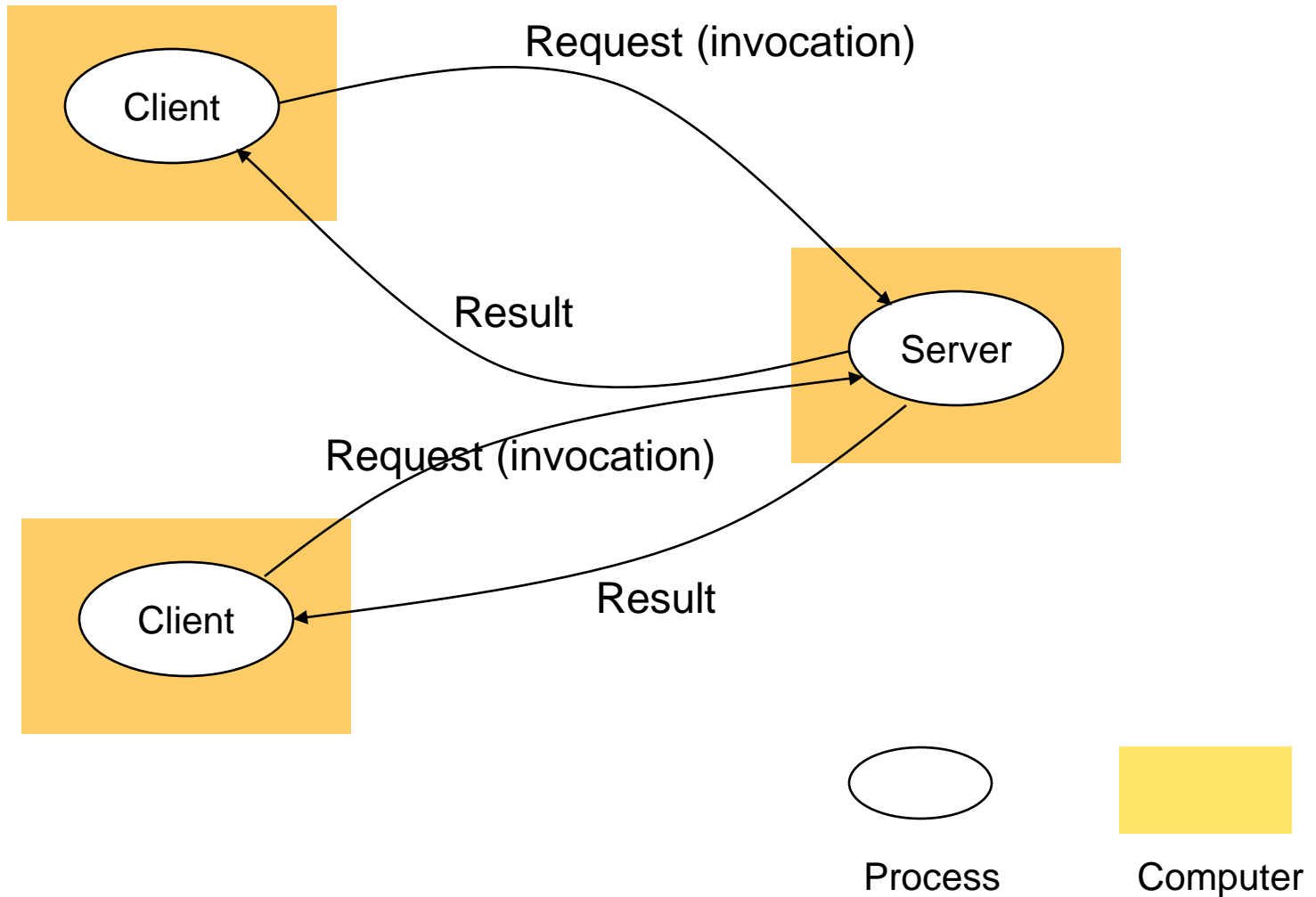
Basic Client–Server Model

Characteristics:

- There are processes offering services (**servers**)
- There are processes that use services (**clients**)
- Clients and servers can be on different machines
- Clients follow request/reply model wrt to using services



Client-Server Communication



Advantages of Client/Server Computing

- More efficient division of labor
- Horizontal and vertical scaling of resources
- Better price/performance on client machines
- Ability to use familiar tools on client machines
- Client access to remote data (via standards)
- Full DBMS functionality provided to client workstations
- Overall better system price/performance

An Example Client and Server (1)

- The *header.h* file used by the client and server.

```
/* Definitions needed by clients and servers.          */
#define TRUE 1
#define MAX_PATH 255 /* maximum length of file name */
#define BUF_SIZE 1024 /* how much data to transfer at once */
#define FILE_SERVER 243 /* file server's network address */

/* Definitions of the allowed operations */
#define CREATE 1 /* create a new file */
#define READ 2 /* read data from a file and return it */
#define WRITE 3 /* write data to a file */
#define DELETE 4 /* delete an existing file */

/* Error codes. */
#define OK 0 /* operation performed correctly */
#define E_BAD_OPCODE -1 /* unknown operation requested */
#define E_BAD_PARAM -2 /* error in a parameter */
#define E_IO -3 /* disk error or other I/O error */

/* Definition of the message format. */
struct message {
    long source; /* sender's identity */
    long dest; /* receiver's identity */
    long opcode; /* requested operation */
    long count; /* number of bytes to transfer */
    long offset; /* position in file to start I/O */
    long result; /* result of the operation */
    char name[MAX_PATH]; /* name of file being operated on */
    char data[BUF_SIZE]; /* data to be read or written */
};
```

An Example Client and Server (2)

- A sample server.

```
#include <header.h>
void main(void) {
    struct message m1, m2;          /* incoming and outgoing messages */
    int r;                          /* result code */

    while(TRUE) {                  /* server runs forever */
        receive(FILE_SERVER, &m1); /* block waiting for a message */
        switch(m1.opcode) {        /* dispatch on type of request */
            case CREATE:           r = do_create(&m1, &m2); break;
            case READ:             r = do_read(&m1, &m2); break;
            case WRITE:            r = do_write(&m1, &m2); break;
            case DELETE:           r = do_delete(&m1, &m2); break;
            default:                r = E_BAD_OPCODE;
        }
        m2.result = r;              /* return result to client */
        send(m1.source, &m2);      /* send reply */
    }
}
```

An Example Client and Server (3)

- A client using the server to copy a file.

```
(a)
#include <header.h>
int copy(char *src, char *dst){
    struct message ml;
    long position;
    long client = 110;

    initialize( );
    position = 0;
    do {
        ml.opcode = READ;
        ml.offset = position;
        ml.count = BUF_SIZE;
        strcpy(&ml.name, src);
        send(FILESERVER, &ml);
        receive(client, &ml);

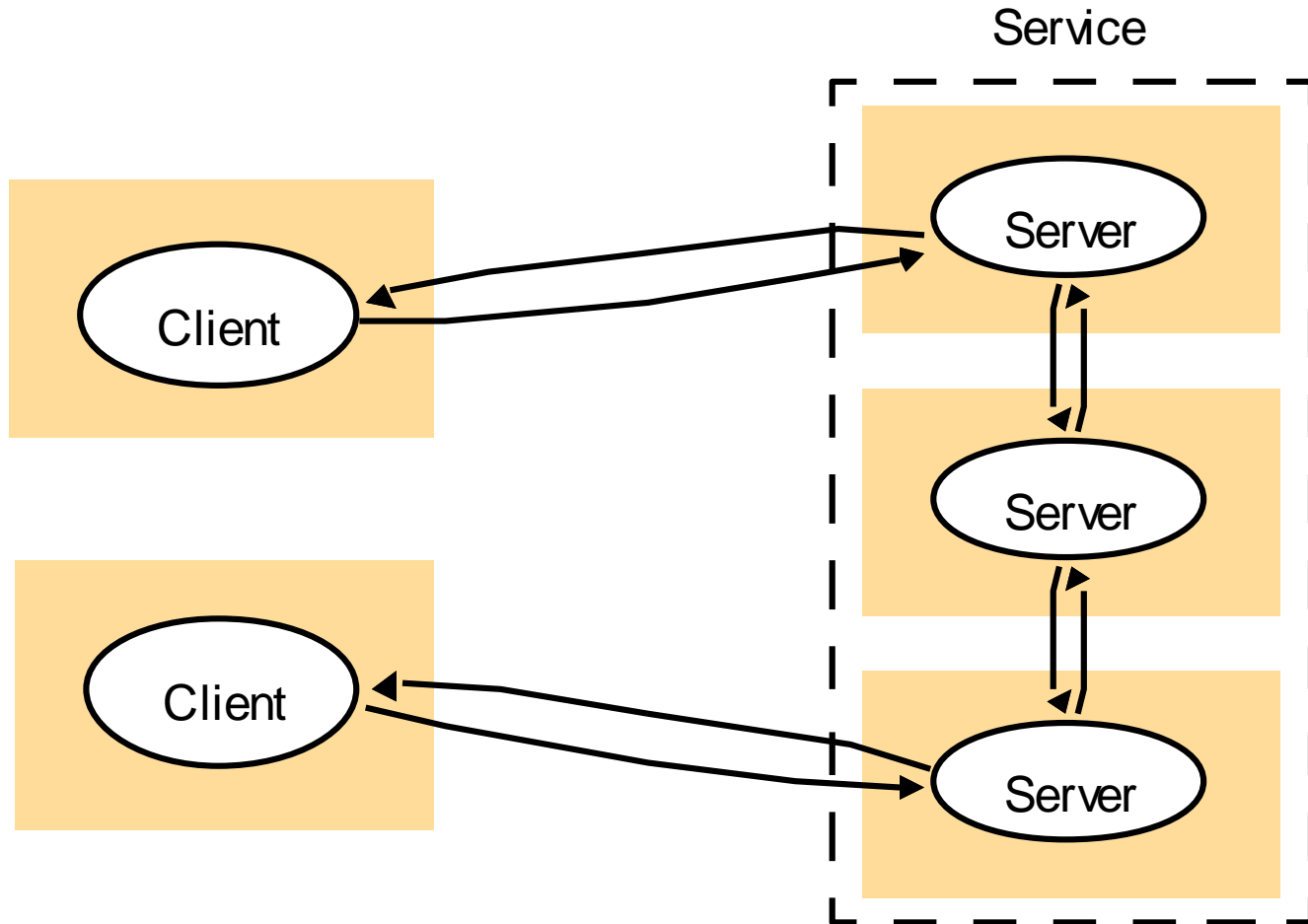
        /* Write the data just received to the destination file.
        ml.opcode = WRITE;
        ml.offset = position;
        ml.count = ml.result;
        strcpy(&ml.name, dst);
        send(FILE_SERVER, &ml);
        receive(client, &ml);
        position += ml.result;
    } while( ml.result > 0 );
    return(ml.result >= 0 ? OK : ml.result);
}
```

/* procedure to copy file using the server */
/* message buffer */
/* current file position */
/* client's address */
/* prepare for execution */
/* operation is a read */
/* current position in the file */
/* how many bytes to read*/
/* copy name of file to be read to message */
/* send the message to the file server */
/* block waiting for the reply */
/* operation is a write */
/* current position in the file */
/* how many bytes to write */
/* copy name of file to be written to buf */
/* send the message to the file server */
/* block waiting for the reply */
/* ml.result is number of bytes written */
/* iterate until done */
/* return OK or error code */

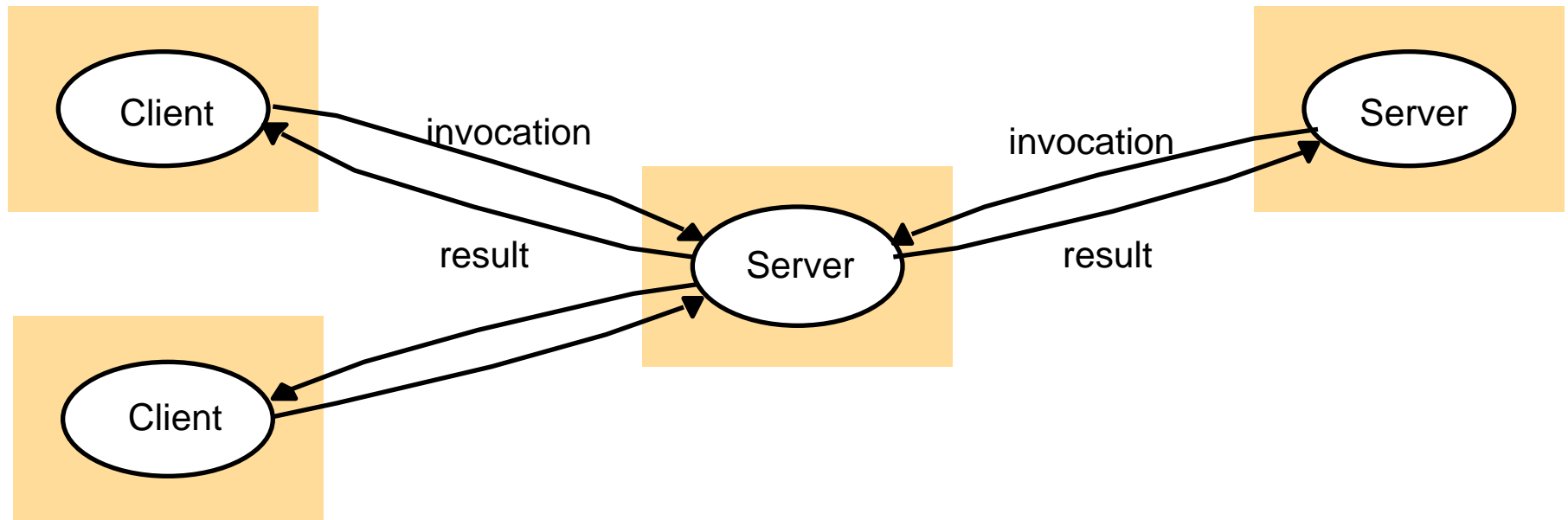
Problems With Multiple-Client/Single Server

- Server forms bottleneck
- Server forms single point of failure
- System scaling difficult

Service Across Multiple Servers



Multiple-Client/Multiple-Server Communication



Application Layering

Traditional three-layered view

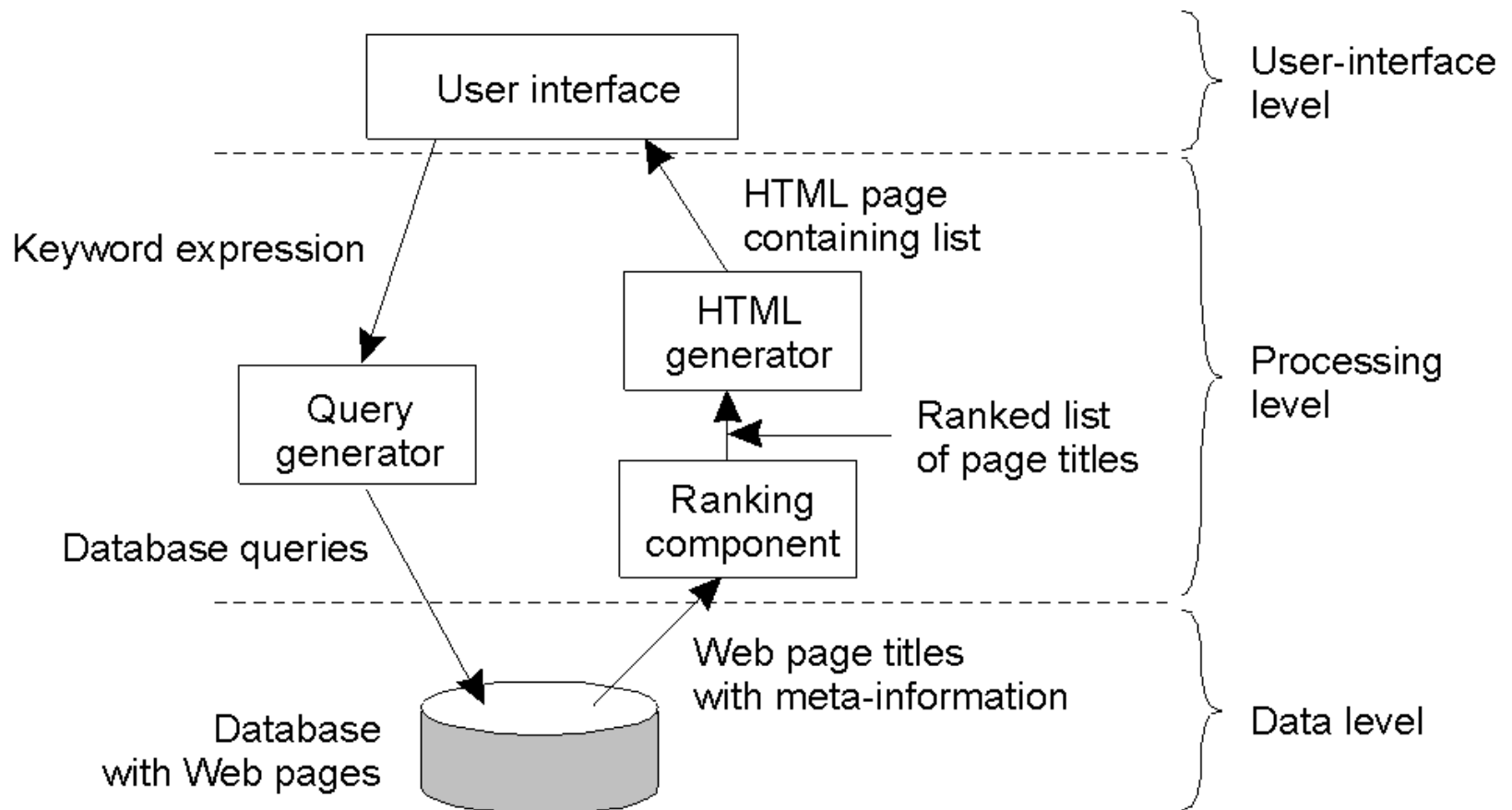
- User-interface layer contains units for an application's user interface
- Processing layer contains the functions of an application, i.e. without specific data
- Data layer contains the data that a client wants to manipulate through the application components

Observation

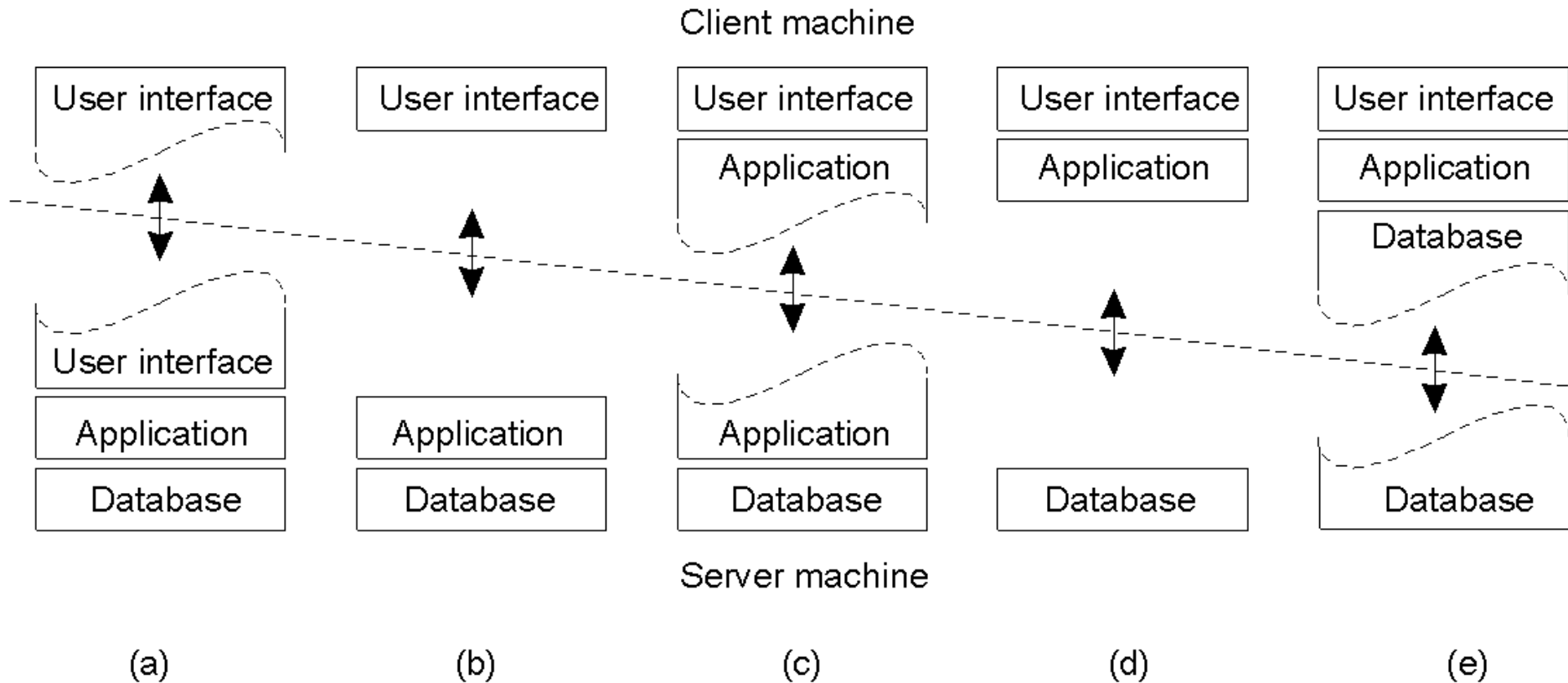
This layering is found in many distributed information systems, using traditional database technology and accompanying applications.

Multitier Systems

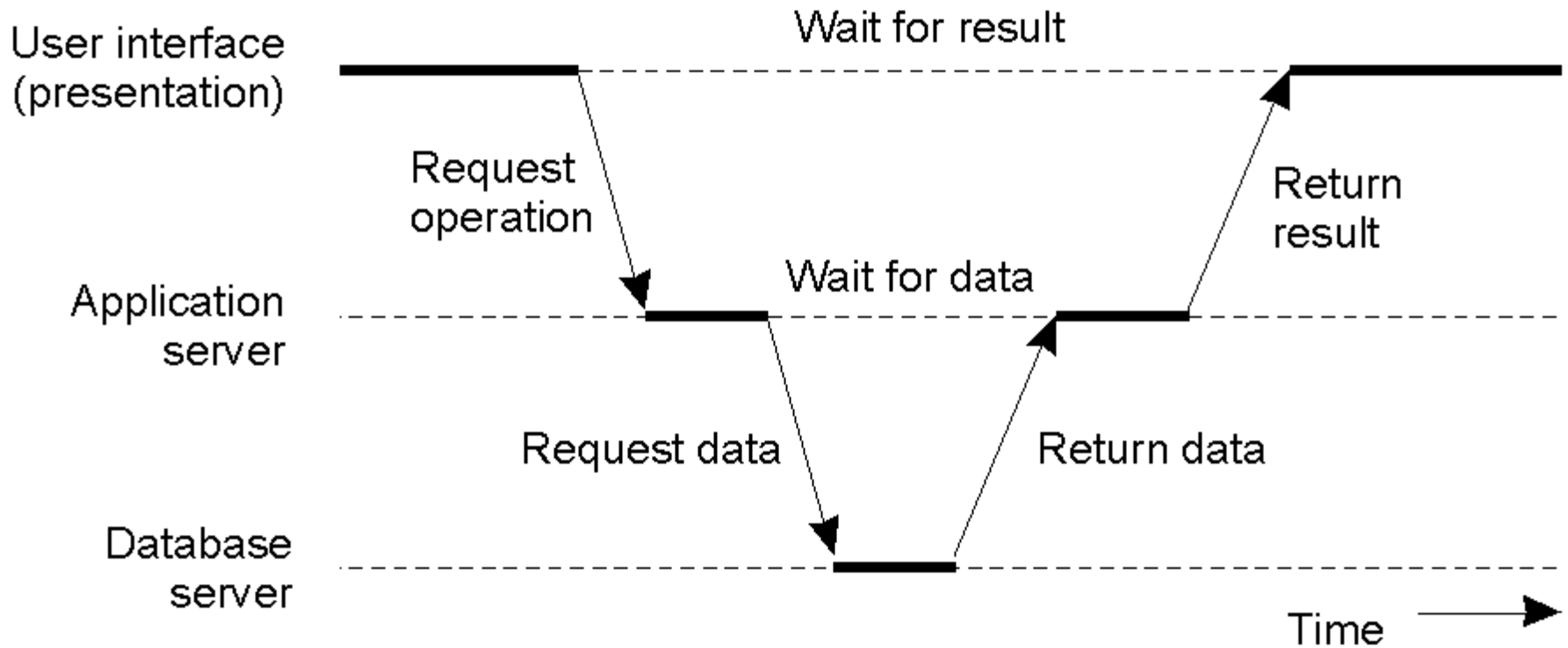
■ Example: Internet Search Engines



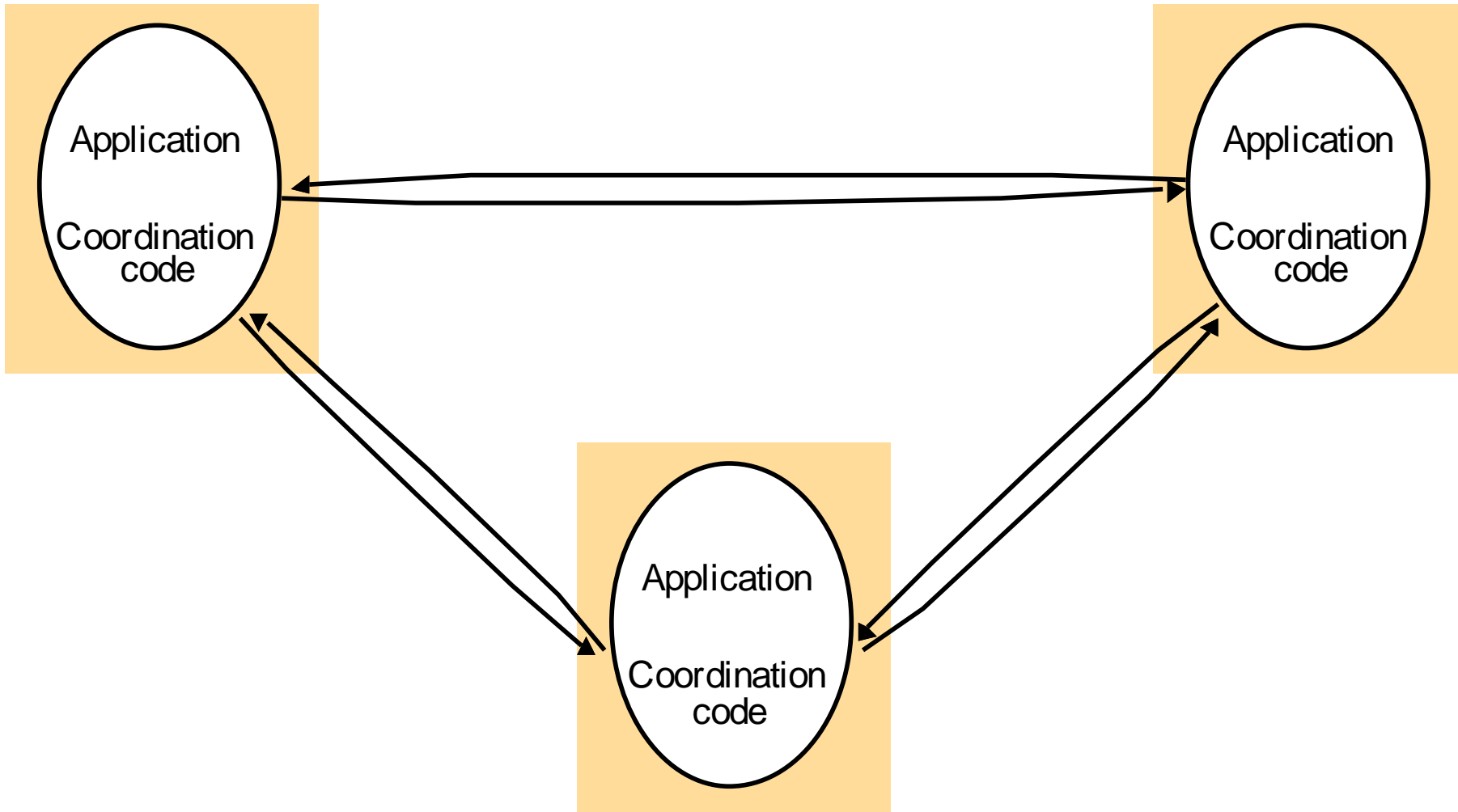
Multitier System Alternatives



Communication in Multitier Systems



Peer-to-Peer Systems



Decentralized Architectures

Observation

In the last couple of years we have been seeing a tremendous growth in **peer-to-peer systems**.

- **Structured P2P**: nodes are organized following a specific distributed data structure
- **Unstructured P2P**: nodes have randomly selected neighbors
- **Hybrid P2P**: some nodes are appointed special functions in a well-organized fashion

Note

In virtually all cases, we are dealing with **overlay networks**: data is routed over connections setup between the nodes (cf. application-level multicasting)

Unstructured P2P Systems

Observation

Many unstructured P2P systems attempt to maintain a **random graph**.

Basic principle

Each node is required to contact a randomly selected other node:

- Let each peer maintain a **partial view** of the network, consisting of c other nodes
- Each node P periodically selects a node Q from its partial view
- P and Q exchange information **and** exchange members from their respective partial views

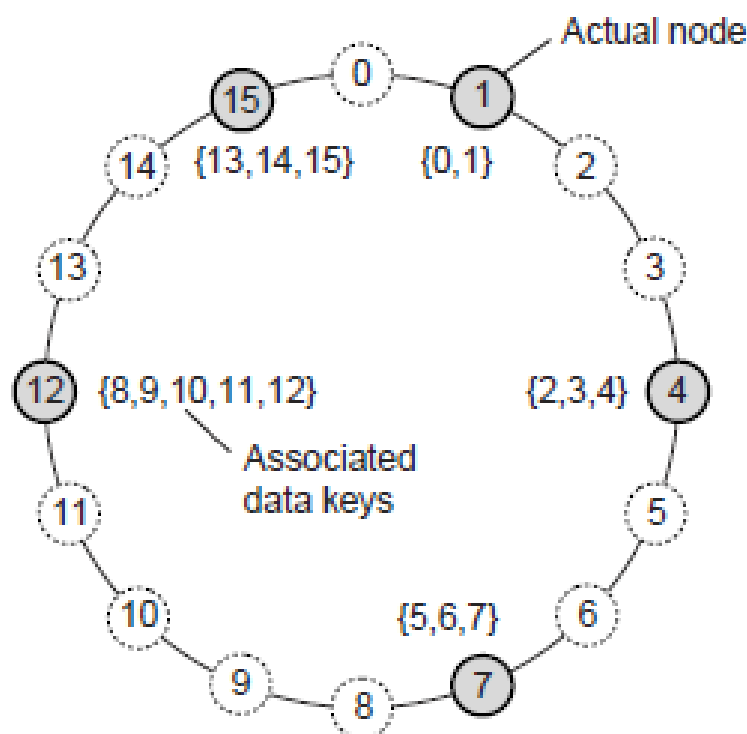
Note

It turns out that, depending on the exchange, randomness, but also **robustness** of the network can be maintained.

Structured P2P Systems

Basic idea

Organize the nodes in a structured **overlay network** such as a logical ring, and make specific nodes responsible for services based only on their ID.



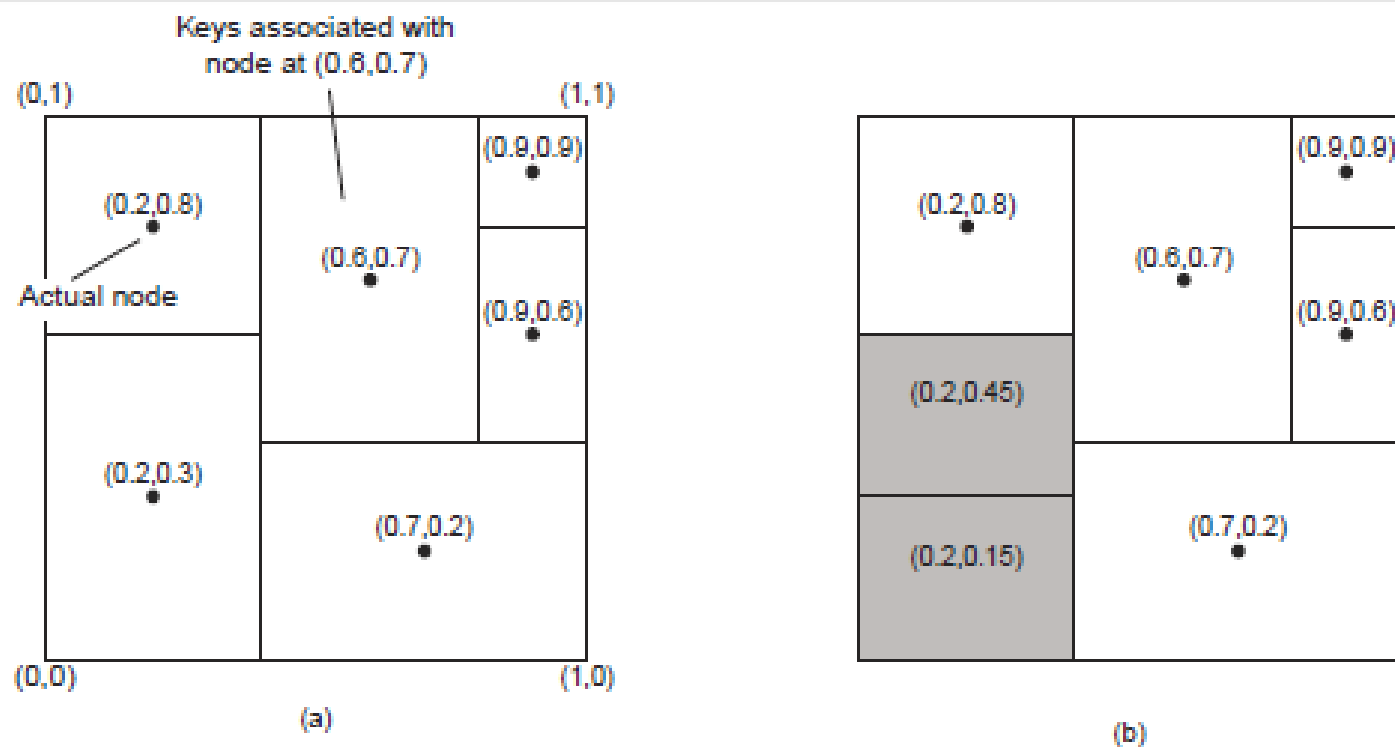
Note

The system provides an operation ***LOOKUP(key)*** that will efficiently **route** the lookup request to the associated node.

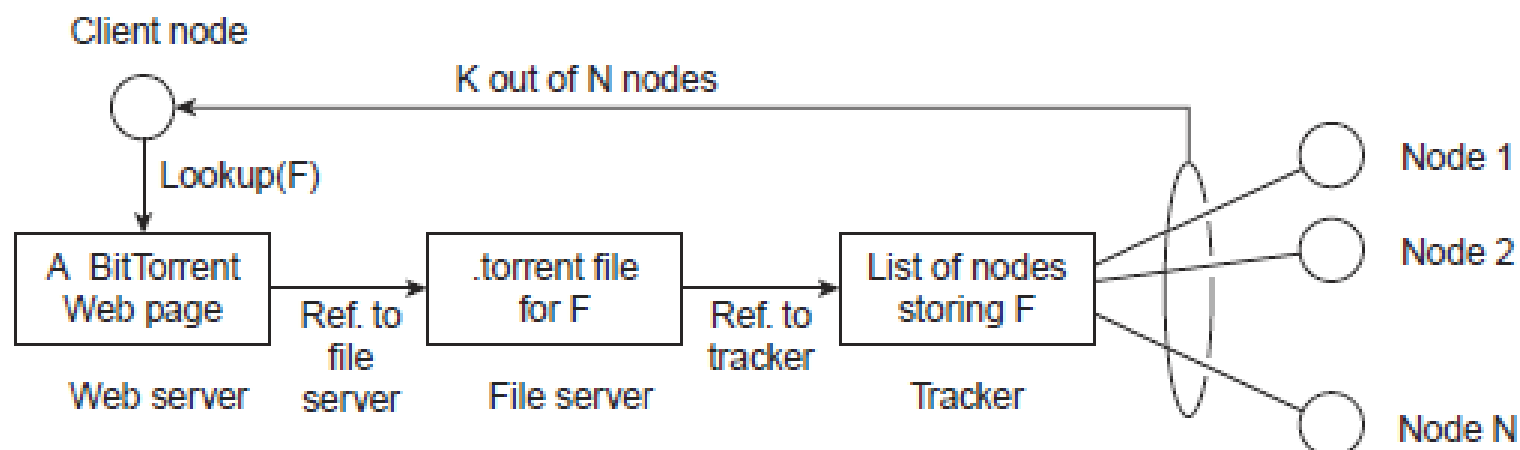
Structured P2P Systems

Other example

Organize nodes in a d -dimensional space and let every node take the responsibility for data in a specific region. When a node joins \Rightarrow split a region.



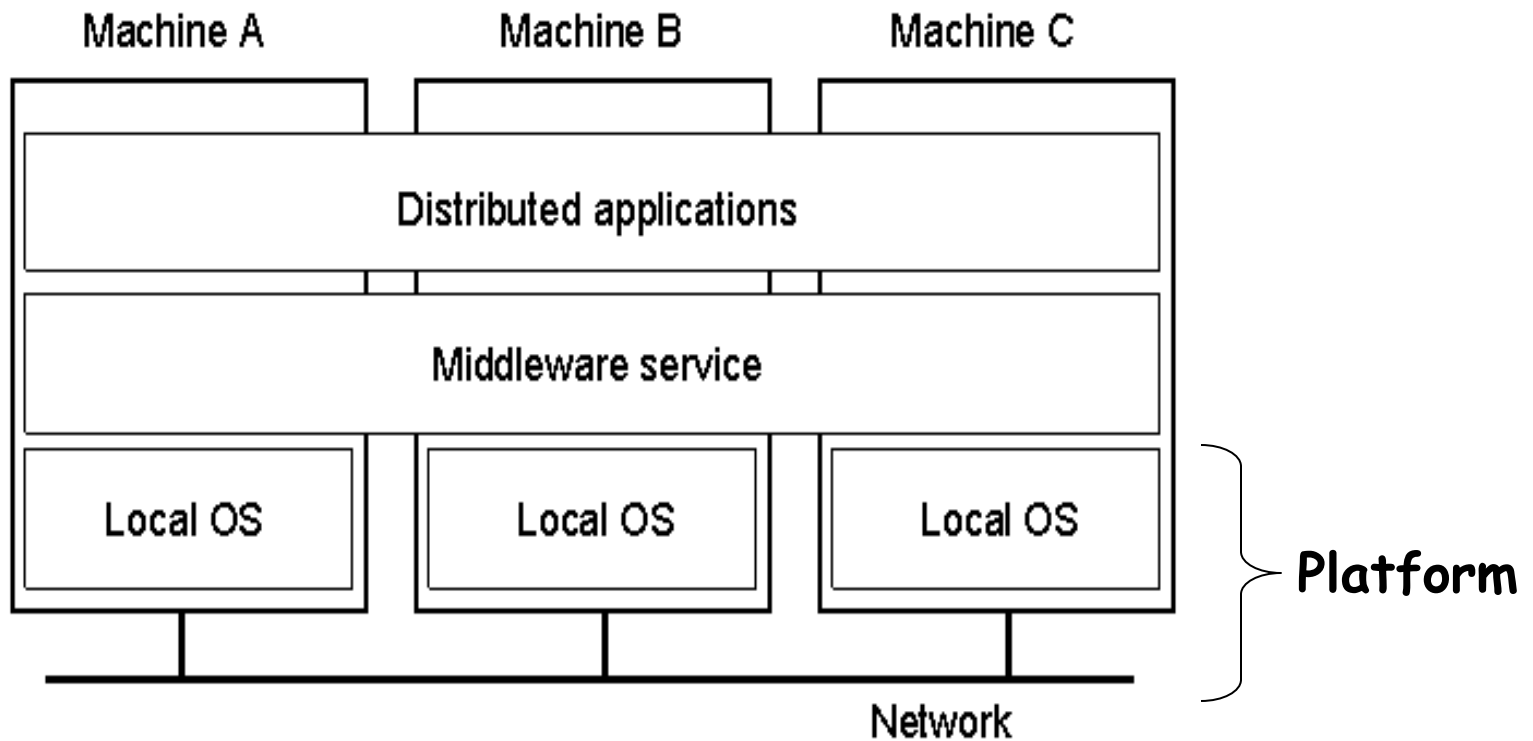
Hybrid Architectures: C/S with P2P – BitTorrent



Basic idea

Once a node has identified where to download a file from, it joins a **swarm** of downloaders who **in parallel** get file chunks from the source, but also distribute these chunks amongst each other.

Software Layers



Layers

- Platform
 - Fundamental communication and resource management services
 - We won't be worried about these
- Middleware
 - Provides a service layer that hides the details and heterogeneity of the underlying platform
 - Provides an “easier” API for the applications and services
 - RPC, RMI, CORBA, etc.
- Applications
 - Distributed applications, services
 - Examples: e-mail, ftp, etc

Example Client/Server Middleware

■ Remote Procedure Call (RPC)

- Uses the well-known procedure call semantics.
- The caller makes a procedure call and then waits. If it is a local procedure call, then it is handled normally; if it is a remote procedure, then it is handled as a remote procedure call.
- Caller semantics is blocked send; callee semantics is blocked receive to get the parameters and a nonblocked send at the end to transmit results.

